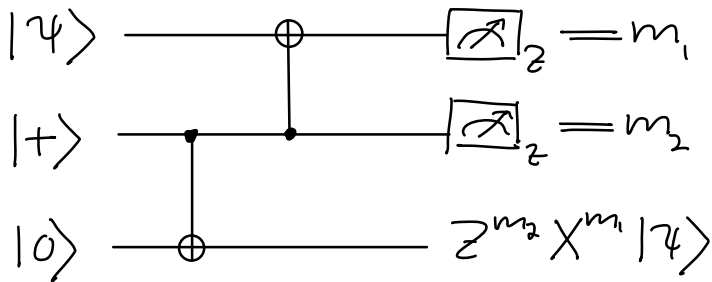
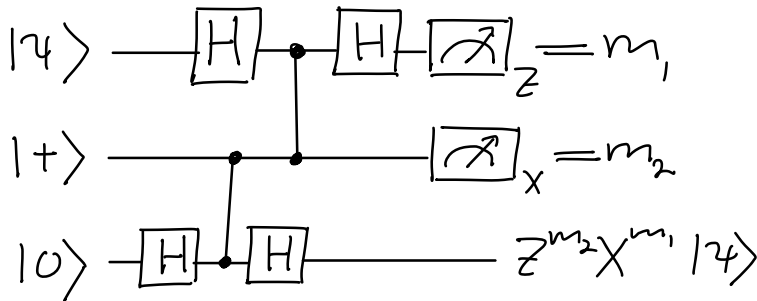


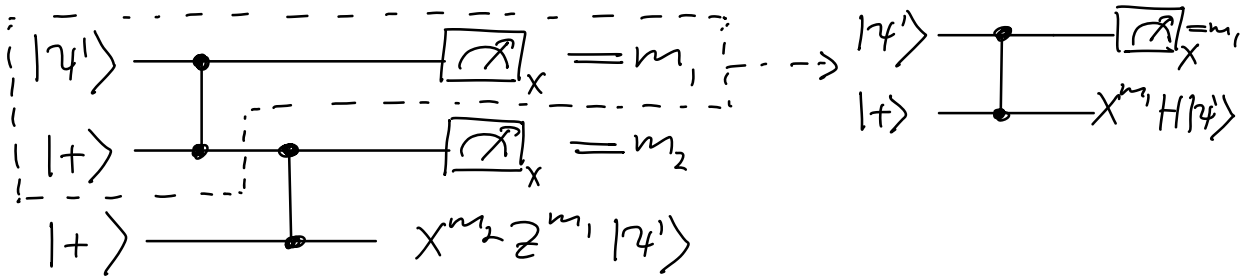
We can decompose the teleportation circuit into two elementary teleportations:



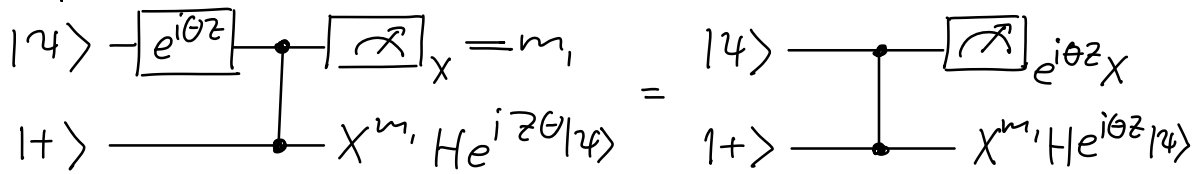
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||



A single-qubit Z rotation $e^{i\theta Z}$ can be implemented in a teleportation-based way:



2) Now we are ready to formulate MBQC

Decompose single-qubit unitary U
(up to a phase) as:

$$U = H e^{i\phi Z} e^{i\theta X} e^{i\gamma Z}$$

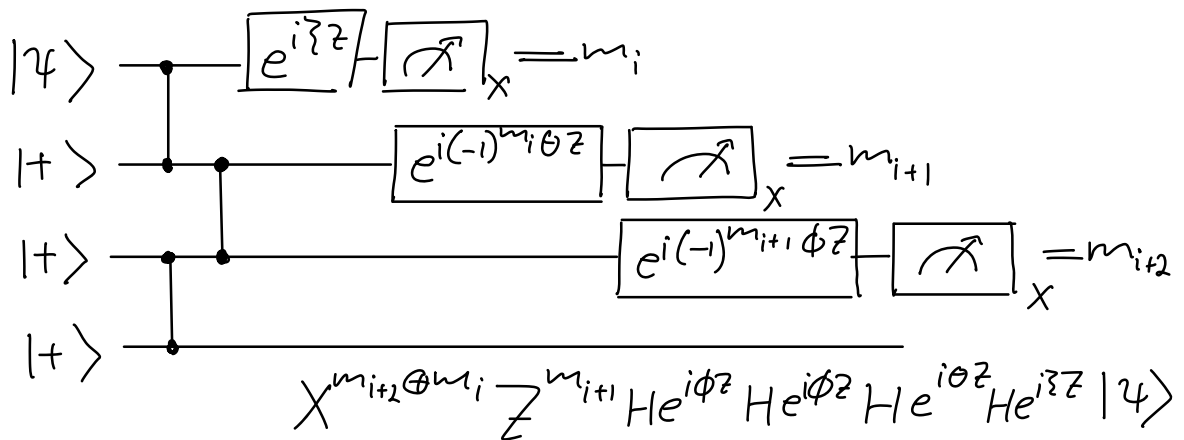
$$= H e^{i\phi Z} H e^{i\theta Z} H e^{i\gamma Z}$$

→ arbitrary single-qubit unitary operation
can be performed by a sequence
of one-bit teleportations!

We have to take care of the byproduct
Pauli operators depending on the
measurement outcomes:

$$U = X^{m_{i+2}} H e^{i\phi Z} X^{m_{i+1}} H e^{i\theta Z} X^{m_i} H e^{i\gamma Z}$$

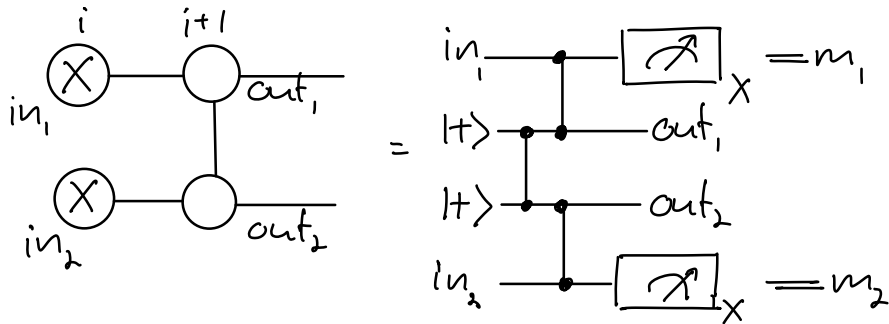
$$= X^{m_{i+2} \oplus m_i} Z^{m_{i+1}} H e^{i(-1)^{m_{i+1}} \phi Z} H e^{i(-1)^{m_{i+1}} \theta Z} H e^{i\gamma Z}$$



By choosing $\theta' = (-1)^{m_i} \theta$ and $\phi' = (-1)^{m_{i+1}} \phi$,
 depending on previous measurement
 outcomes, U can be forwarded
 → "feed forward"

3) Measurement-based two-qubit gate:

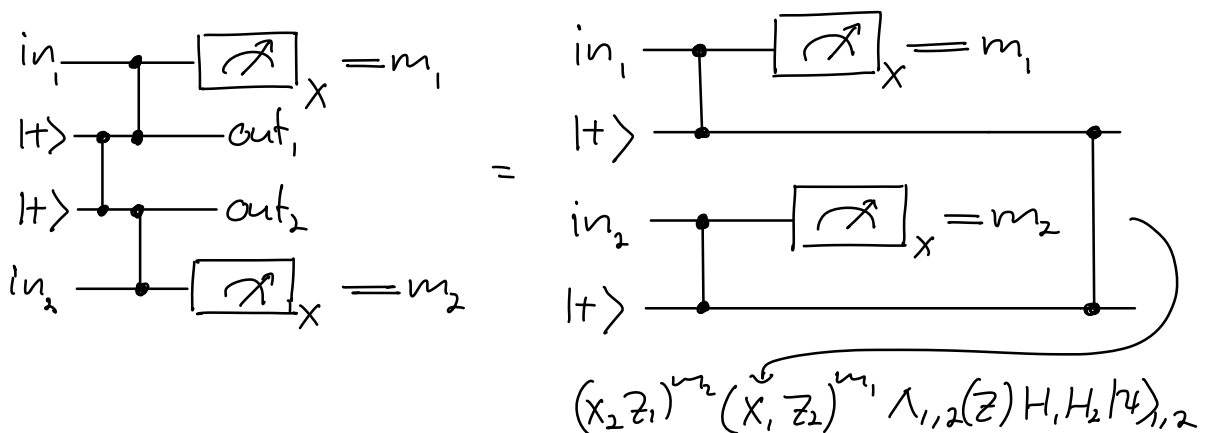
resource state is cluster state:



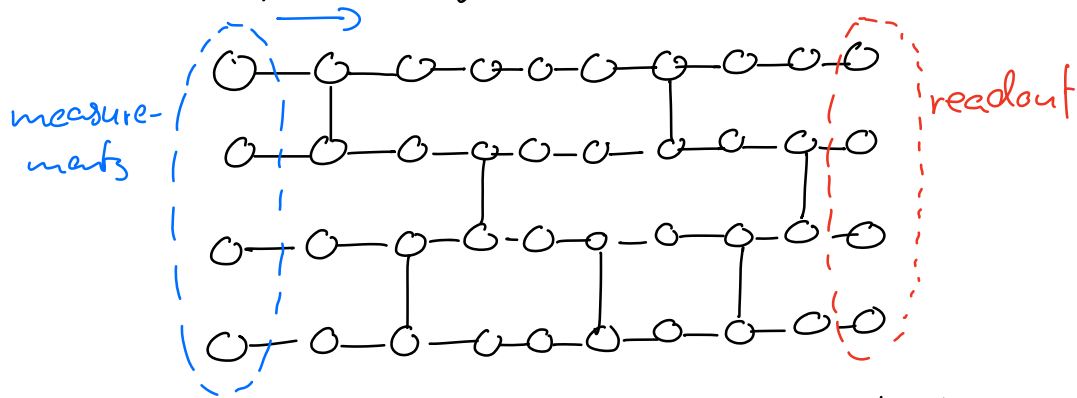
(recall graph state

$$|G\rangle = \prod_{(i,j) \in E} \Lambda(z)_{i,j} (|+\rangle^{\otimes |V|})$$

Now observe that



More generally, one can prepare an n -qubit graph-state :



→ perform measurements starting from the left-most column and successively to the right

→ simulate universal quantum computation by measurements!

§ 2.7 Quantum Error Correction Codes

i) Three-Qubit Bit-Flip Code :

Consider stabilizer generators

$$S_1 = Z_1 Z_2, \quad S_2 = Z_2 Z_3$$

→ stabilizer subspace is spanned by

$$|0_L\rangle = |000\rangle, \quad |1_L\rangle = |111\rangle$$

define

- logical Pauli-X operator $L_X \equiv X_1 X_2 X_3$
- logical Pauli-Z operator $L_Z \equiv Z_1$

Consider a bit flip error with an error probability p :

$$\mathcal{E}_i \rho = (1-p)\rho + pX_i \rho X_i$$

acting successively on initial state

$$|\psi_L\rangle = \alpha |0_L\rangle + \beta |1_L\rangle$$

$$\rightarrow \mathcal{E}_1 \circ \mathcal{E}_2 \circ \mathcal{E}_3 |\psi_L\rangle \langle \psi_L|$$

$$= (1-p)^3 |\psi_L\rangle \langle \psi_L| + p(1-p)^2 \sum_i X_i |\psi_L\rangle \langle \psi_L| X_i + O(p^2)$$

error X_i maps code to an orthogonal subspace

→ perform projective measurements onto orthogonal subspaces,

$$P_k^\pm = (I \pm S_k)/2$$

→ measurement result is called "error syndrome"

4 error syndromes:

$$P_0 \equiv |000\rangle\langle 000| + |111\rangle\langle 111| \quad \text{no error}$$

$$P_1 \equiv |100\rangle\langle 100| + |011\rangle\langle 011| \quad \text{bit flip on qubit one}$$

$$P_2 \equiv |010\rangle\langle 010| + |101\rangle\langle 101| \quad \text{bit flip on qubit two}$$

$$P_3 \equiv |001\rangle\langle 001| + |110\rangle\langle 110| \quad \text{bit flip on qubit three}$$

suppose corrupted state is $a|100\rangle + b|011\rangle$

$$\rightarrow \langle \mathcal{U} | P_1 | \mathcal{U} \rangle = 1$$

Note: the syndrome measurement preserves the state $\alpha|0\rangle + \beta|1\rangle$

Recovery:

$$R \circ \mathcal{E}_1 \circ \mathcal{E}_2 \circ \mathcal{E}_3 |\psi_L\rangle \langle \psi_L|$$

$$= [(1-p)^3 + 3p(1-p)^2] |\psi_L\rangle \langle \psi_L| + \mathcal{O}(p^2)$$

where the recovery operator is given by

$$R\rho = P_1^+ P_2^+ \rho P_2^+ P_1^+ + X_1 P_1^- P_2^+ \rho P_2^+ P_1^- X_1$$

$$+ X_2 P_1^- P_2^- \rho P_2^- P_1^- X_2 + X_3 P_1^+ P_2^- \rho P_2^- P_1^+ X_3$$

"fidelity" between a pure and a mixed state is given by $\sqrt{\langle \psi | \rho | \psi \rangle}$ without error correction:

$$\rho = (1-p) |\psi\rangle \langle \psi| + p X |\psi\rangle \langle \psi| X$$

$$\Rightarrow F = \sqrt{\langle \psi | \rho | \psi \rangle} = \sqrt{(1-p) + p \underbrace{\langle \psi | X | \psi \rangle \langle \psi | X | \psi \rangle}_{\geq 0}}$$

and = 0 for $|\psi\rangle = 0$

$$\Rightarrow F \geq \sqrt{1-p}$$

with error correction we get

$$F = \sqrt{\langle \psi | \rho | \psi \rangle} \geq \sqrt{(1-p)^3 + 3p(1-p)^2}$$

\rightarrow improved for $p < \frac{1}{2}$!